

Optimal Partitioning Of Multi-Block Structured Grids

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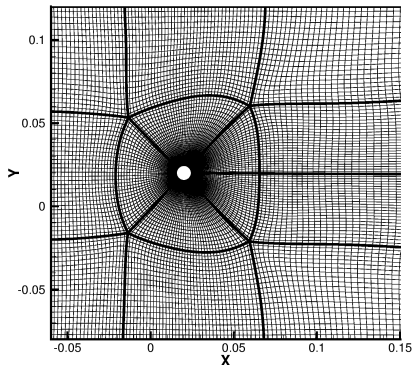
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- The Advanced Aerodynamics Department uses the FANSC software to solve the Navier-Stokes equations.
- These are coupled non-linear partial differential equations that govern the motion of compressible, viscous fluids such as air.
- The Navier-Stokes equations are solved on a fluid volume region discretized into a large number of cells.
- The solution of the equations yields the values of the density, velocity, and pressure within each cell.
- The computation is distributed over many processors (1000 today, maybe 10000 in a few years from now...).

Multi-block Structured Grids

The discretization produces regions that “look like” grids.



The union of several regions may (or may not) look like a grid.

Allocating Computations to Processors: the Objectives

One must find a way of allocating groups of cells to processors. A group of cells will consist of several grid-like regions called “blocks”.

- The computational load must be balanced, i.e., the difference between the numbers of cells allocated to any two processors must be as small as possible.
- The maximum communication cost between a processor and all the other processors must be as small as possible.
- The number of blocks allocated to a given processor must not be too large.
- None of the blocks allocated to a given processor must be too small or too “thin.”

Decomposition of the Problem

Since the problem is too complicated to be solved as a whole, we have decomposed it into subproblems.

- ① Refine the original partition (produced by an engineer) so as to obtain small blocks and create a graph with weights on nodes and edges.
- ② Partition the graph into clusters (i.e., groups of blocks assigned to a given processor).
- ③ Reorganize the blocks within a given cluster so as to obtain relatively large blocks that are grid-like and are not too “thin.”

Topology Refinement and Graph Structure

```
# BLOCKS
5
=====Block 1=====
Imax 65 Jmax 65 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 1 65 97 0 2 0
SYM 97 1 1 1 161 97 0
SYM 1 1 1 1 65 97 0
SYM 65 1 1 1 65 97 0
WAL 1 1 1 1 65 97 0
FAR 1 1 97 65 97 0

=====Block 2=====
Imax 257 Jmax 65 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 97 1 1 1 161 97 0
CON 1 1 65 1 65 97 0
CON 1 1 1 1 65 97 0
CON 2 257 1 1 257 65 97 0
CON 257 1 1 257 65 97 0
CON 2 1 1 1 65 97 0
CON 1 1 1 33 1 97 0
CON 5 33 1 1 1 97 0
CON 225 1 1 257 1 97 0
CON 5 65 1 1 33 1 97 0
CON 1 65 1 257 65 97 0
CON 3 1 1 257 1 97 0
SYM 33 1 1 97 1 97 0
SYM 161 1 1 225 1 97 0
WAL 1 1 1 257 65 1 100
FAR 1 1 97 257 65 97 0

=====Block 3=====
Imax 257 Jmax 33 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 257 1 97 0
CON 2 1 65 1 257 65 97 0
CON 1 1 1 1 33 97 0
CON 3 257 1 1 257 33 97 0
CON 1 1 1 257 33 97 0
CON 4 1 33 1 257 33 97 0
CON 1 1 1 257 1 97 0
WAL 1 1 1 257 33 1 1
FAR 1 1 97 257 33 97 0

=====Block 4=====
Imax 257 Jmax 65 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 257 1 97 0
CON 3 1 33 1 257 33 97 0
CON 1 1 1 1 65 97 0
CON 4 257 1 1 257 65 97 0
CON 257 1 1 257 65 97 0
CON 4 1 1 1 65 97 0
CON 1 65 1 129 65 97 0
CON 4 257 65 1 129 65 97 0
CON 129 65 1 257 65 97 0
CON 4 129 65 1 65 97 0
WAL 1 1 1 257 65 1 1
FAR 1 1 97 257 65 97 0

=====Block 5=====
Imax 65 Jmax 129 Kmax 97 Total surfaces=====
7
```

Figure 1: Topology file

Topology Refinement and Graph Structure

```
# BLOCKS
5
=====Block 1=====
Imax 65 Jmax 65 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 1 65 97 0 2 0
SYM 2 97 1 1 161 97 0
SYM 1 1 1 1 65 97 0
SYM 65 1 1 65 65 97 0
WAL 1 1 1 1 65 1 97 0
FAR 1 1 1 97 65 97 100
=====Block 2=====
Imax 257 Jmax 65 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 97 1 1 1 161 97 0
CON 1 1 65 1 65 97 0 2 0
CON 1 1 1 1 65 97 0
CON 2 257 1 1 257 65 97 0 1 2
CON 257 1 1 257 65 97 0
CON 2 1 1 1 65 97 0 1 2
CON 1 1 1 33 1 97 0
CON 5 33 1 1 1 97 0 2 0
CON 225 1 1 257 1 97 0
CON 5 65 1 1 33 1 97 0 2 0
CON 1 1 65 1 257 65 97 0
CON 3 1 1 1 257 1 97 0 2 0
SYM 33 1 1 97 1 97 0
SYM 161 1 1 225 1 97 0
WAL 1 1 1 257 65 1 100
FAR 1 1 1 97 257 65 97 0
=====Block 3=====
Imax 257 Jmax 33 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 1 257 1 97 0
CON 2 1 65 1 257 65 97 0 2 0
CON 1 1 1 1 33 97 0
CON 3 257 1 1 257 33 97 0 1 2
CON 1 257 1 1 257 33 97 0
CON 4 1 1 1 33 97 0 1 2
CON 1 33 1 257 33 97 0
CON 4 1 1 1 257 1 97 0 2 0
WAL 1 1 1 257 33 1 1
FAR 1 1 1 97 257 33 97 0
=====Block 4=====
Imax 257 Jmax 65 Kmax 97 Total surfaces=====
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 1 257 1 97 0
CON 3 1 33 1 257 33 97 0 2 0
CON 1 1 1 1 65 97 0
CON 4 257 1 1 257 65 97 0 1 2
CON 257 1 1 257 65 97 0
CON 4 1 1 1 65 97 0 1 2
CON 1 65 1 129 65 97 0
CON 4 257 65 1 129 65 97 0 2 0
CON 129 65 1 257 65 97 0
CON 4 129 65 1 65 97 0 2 0
WAL 1 1 1 257 65 1 1
FAR 1 1 1 97 257 65 97 0
=====Block 5=====
Imax 65 Jmax 129 Kmax 97 Total surfaces=====
```



Figure 1: Topology file

Topology Refinement and Graph Structure

```
# BLOCKS
5
=====Block 1=====
Imax Jmax Kmax Total surfaces=====
1 65 65 97 6
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 161 1 97 0 2 0
SYM 97 1 1 1 65 97 0
SYM 65 1 1 65 65 97 0
SYM 1 1 1 65 1 97 0
WAL 1 1 1 65 65 1 100
FAR 1 1 97 65 97 0
=====Block 2=====
Imax Jmax Kmax Total surfaces=====
2 257 65 97 10
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 97 1 1 161 1 97 0 2 0
CON 1 1 65 1 65 97 0
CON 1 1 1 1 65 97 0
2 257 1 1 257 65 97 0 1 2
CON 257 1 1 257 65 97 0
2 1 1 1 65 97 0 1 2
CON 1 1 1 33 1 97 0 2 0
5 33 1 1 1 97 0
CON 225 1 1 257 1 97 0 2 0
5 65 1 1 33 1 97 0
CON 1 65 1 257 65 97 0
3 1 1 1 257 1 97 0 2 0
SYM 33 1 1 97 1 97 0
SYM 161 1 1 225 1 97 0
WAL 1 1 1 257 65 1 100
FAR 1 1 97 257 65 97 0
=====Block 3=====
Imax Jmax Kmax Total surfaces=====
3 257 33 97 6
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 257 1 97 0 2 0
CON 2 1 65 1 257 65 97 0
CON 1 1 1 1 33 97 0 1 2
CON 257 1 1 257 33 97 0
3 1 1 1 33 97 0 1 2
CON 1 33 1 257 33 97 0
4 1 1 1 257 1 97 0 2 0
WAL 1 1 1 257 33 1 1
FAR 1 1 97 257 33 97 0
=====Block 4=====
Imax Jmax Kmax Total surfaces=====
4 257 65 97 7
---Bctype Ista Jsta Ksta Iend Jend Kend icomp iiconnect---
CON 1 1 1 257 1 97 0 2 0
3 1 33 1 257 33 97 0
CON 1 1 1 1 65 97 0 1 2
CON 257 1 1 257 65 97 0
4 1 1 1 65 97 0 1 2
CON 1 65 1 129 65 97 0 2 0
CON 4 257 65 1 129 65 97 0
CON 129 65 1 257 65 97 0 2 0
4 129 65 1 65 97 0
WAL 1 1 1 257 65 1 1
FAR 1 1 97 257 65 97 0
=====Block 5=====
Imax Jmax Kmax Total surfaces=====
5 65 129 97 7
```

Figure 1: Topology file

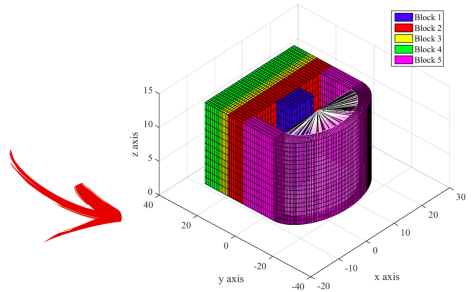


Figure 2: Three dimensional plot of the surface provided

Grid Examples Used For The Partitioning Model

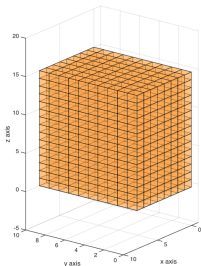


Figure 3: $8 \times 10 \times 15$ structured grid

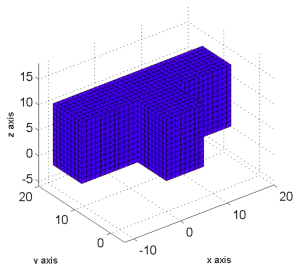


Figure 4: T shape object formed by 2 structured grids

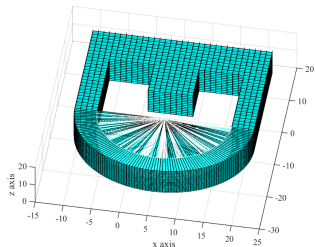


Figure 5: Shape object formed by 3 structured grids

Partitioning Problem Of The Block Graphs (1)

Idea: Turn our problem into a balanced graph clustering problem.

Input: Graph $G = (V, E)$, node weights $w_n : V \mapsto \mathbb{N}_+$, edge weights $w_e : E \mapsto \mathbb{N}_+$, number of colours k .

Objective: Find a colouring of nodes $c : V \mapsto P$ where $P = \{1, \dots, k\}$ such that $T = c_1 \max_{p \in P} V_p + c_2 \max_{p \in P} B_p$ is minimized. Here the volume of a partition is $V_p = \sum_{u \in V} w_n(u) \delta(c(u), p)$ and the boundary of a partition is $B_p = \sum_{(u,v) \in E} w_e((u,v)) [1 - \delta(c(u), c(v))] \delta(c(u), p)$, where δ denotes the indicator function.

Methods for balanced graph clustering:

- Heuristic based on label propagation
- Integer programming formulation
- Freely available programs: KaHIP and METIS (slightly different objective functions)

KaHIP Applied To Structured Grids

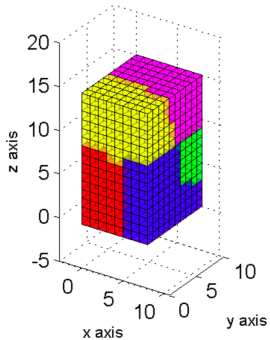
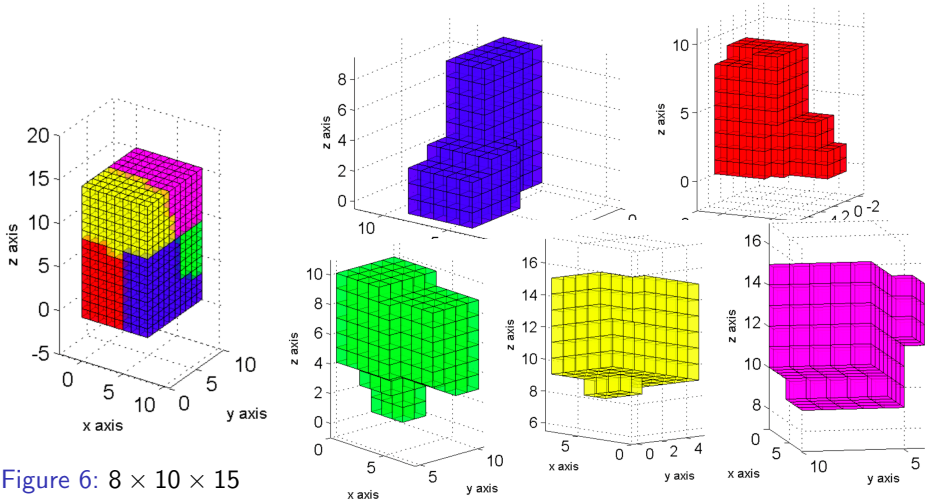


Figure 6: $8 \times 10 \times 15$
structured grid

KaHIP Applied To Structured Grids



Heuristic Partitioning Method Applied To Structured Grid

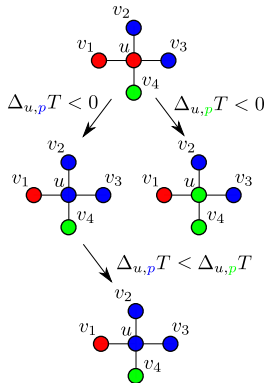
0. Initialize

- Set nodes to random partitions
- Use previous partition as a starting point

1. For each node $u \in V$ (in random order):

- Find colours in the neighborhood of u : $C_{\text{candidates}} = \{c(v) : (u, v) \in E\}$
- For each $p \in C_{\text{candidates}}$ find change in the objective function $\Delta_{u,p}T$ if $c(u)$ is set to p .
- Select p^* that minimizes $\Delta_{u,p}T$. If $\Delta_{u,p^*}T < 0$, set $c(u)$ to p^* .

2. Repeat point 1, if at least one colour was changed in last run of point 1.



Heuristic Partitioning Method Applied To Structured Grid

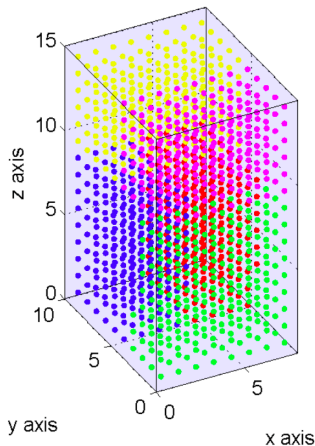


Figure 7: $8 \times 10 \times 15$ Structured Grid

Heuristic Partitioning Method Applied To Structured Grids

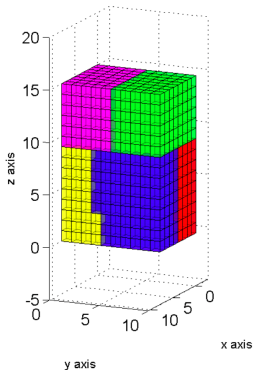


Figure 8: $8 \times 10 \times 15$
structured grid
partitioned by our
algorithm

Heuristic Partitioning Method Applied To Structured Grids

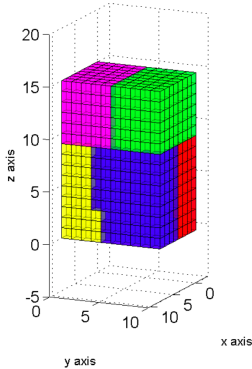


Figure 8: $8 \times 10 \times 15$ structured grid partitioned by our algorithm

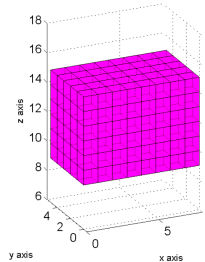
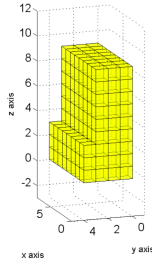
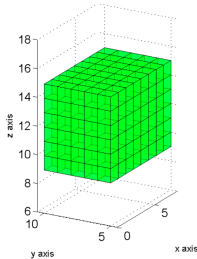
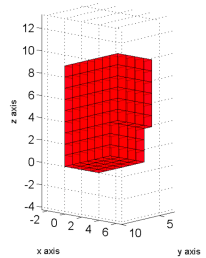
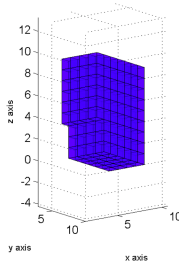
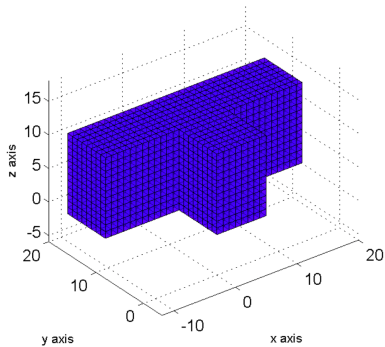


Figure 9: Partitions assigned to 5 processors

Heuristic Partitioning Method Applied To T-shape Grid



Heuristic Partitioning Method Applied To T-shape Grid

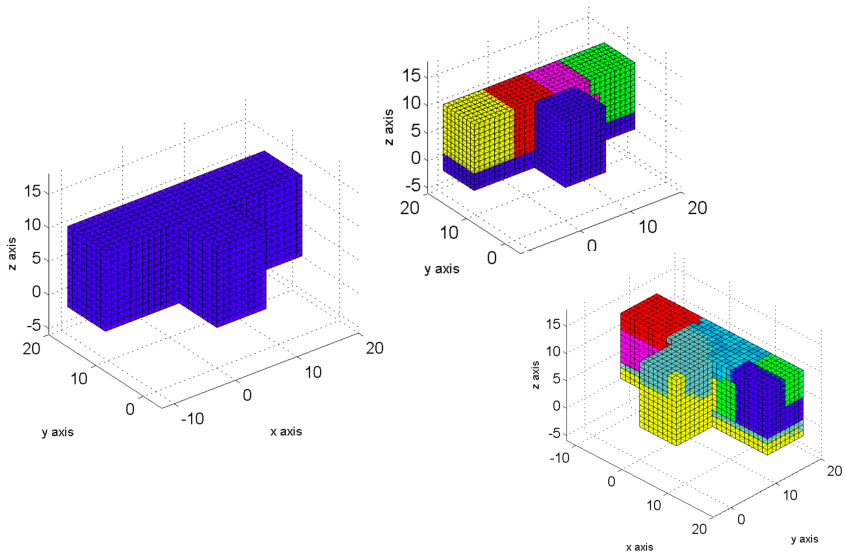


Figure 10: Structure partitioned to 5 and to 8 processors

Partitioning Problem Of The Block Graphs (2)

- Assign the work units to processors in order to minimize the sum of the maximum load and the maximum communication cost
- Mathematical formulation

minimize

$$\left(C_1 \max_{p \in P} \left\{ \sum_{v \in V} w_v y_{vp} \right\} + C_2 \max_p \left\{ \sum_{p^* \neq p} \sum_{(v,u) \in E} w_{(v,u)} y_{vp} y_{up^*} \right\} \right)$$

such that

- ① $\sum_{p \in P} y_{vp} = 1, \forall v \in V$
- ② $y_{vp} \in \{0, 1\}, \forall p \in P, v \in V$

Partitioning Problem Of The Block Graphs (2)

Here is an equivalent formulation.

$$\text{minimize } (C_1 T_1 + C_2 T_2)$$

such that

- ① $\sum_{v \in V} w_v y_{vp} \leq T_1, \forall p \in P$
- ② $\sum_{p^* \neq p} \sum_{(v,u) \in E} w_{(v,u)} y_{vp} y_{up^*} \leq T_2, \forall p \in P$
- ③ $\sum_{p \in P} y_{vp} = 1, \forall v \in V$
- ④ $y_{vp} \in \{0, 1\}, \forall p \in P, v \in V$

Partitioning Problem Of The Block Graphs (2)

We now linearize the term $y_{vp}y_{up^*}$.

1

$$x_{vu}^{pp^*} = \begin{cases} 1 & \text{if node } v \text{ (resp. } u \text{) belongs to cluster } p \text{ (resp. } p^*) \\ 0 & \text{otherwise} \end{cases}$$

2

$$\begin{cases} x_{vu}^{pp^*} \geq y_{vp} + y_{up^*} - 1 \\ x_{vu}^{pp^*} \leq y_{vp} \\ x_{vu}^{pp^*} \leq y_{up^*} \end{cases} \quad \forall (u, v) \in E; p, p^* \in P : p^* \neq p$$

- 3 The minimization structure of the model implies that x will automatically be binary and the last two constraints are not necessary. We redefine the variables.

$$x_e^p = \begin{cases} 1 & \text{if } e \text{ is an outgoing arc from cluster } p \\ 0 & \text{otherwise} \end{cases}$$

- 4 $x_{(v,u)}^p \geq y_{vp} + y_{up^*} - 1, \forall (u, v) \in E; p, p^* \in P, p^* \neq p$
- 5 $x_{(v,u)}^p \geq 0, \forall (u, v) \in E, p \in P$

Partitioning Problem Of The Block Graphs (2)

Here is the final formulation, assuming that all vertex weights are identical and all edge weights are also identical.

$$\text{minimize } C_1 w_1 T_1 + C_2 w_2 T_2$$

such that

- 1 $\sum_{v \in V} y_{vp} \leq T_1, \forall p \in P$ and $\sum_{p \in P} y_{vp} = 1, \forall v \in V$
- 2 $\sum_{e \in E} x_e^p \leq T_2, \forall p \in P$
- 3 $x_{(v,u)}^p \geq y_{vp} + y_{up^*} - 1, \forall (u, v) \in E; p, p^* \in P, p^* \neq p$
- 4 $y_{vp} \in \{0, 1\}, \forall p \in P, v \in V$
- 5 $x_{(v,u)}^p \geq 0, \forall (v, u) \in E; p \in P$

Partitioning Problem Of The Block Graphs (2)

Conclusions after implementing our algorithm.

- Our problem is a very challenging combinatorial problem.
 - It is highly degenerate.
 - The LP is very weak (the two halves of each node are in different clusters).
 - The network size is massive.
- We tried the following approaches.
 - Cplex
 - Benders decomposition
 - MIP-based heuristic (I, II)

If we know the values of the y variables, the communication cost is almost trivial to compute.

Pseudocode of MIP-based heuristic (II)

- ① Find the degree of each node.
- ② For each cluster (associated to a given processor), execute the following steps.
 - ① Find the node with the smallest degree (i.e., v).
 - ② While the load limit is not reached and some neighbour exists, do
 - ① Find all neighbours of v and add them to the current list;
 - ② Update v by selecting one of its neighbours for exploration.
 - ③ Fix the binary variables associated with each node in the cluster list.
- ③ Solve the “restricted MIP.”

Partitioning Problem Of The Block Graphs (2)

Results (the time limit was set to 30 minutes)

Approach	Instance I	Instance II
Cplex	149951	1966080
Benders	150480	N/A
Heuristic I	162240	369280
Heuristic II	130880	291136

Reorganizing the Blocks Assigned to a Given Processor

We have considered two approaches.

- ① Treat the union of all the blocks as a single structure and try to decompose it into “cuboids”, i.e., grid-like regions.
- ② Attempt to merge the blocks themselves by applying some algorithm (probably a greedy algorithm).

The first approach requires very complicated and time-consuming algorithms and the second one does not guarantee that the solution produced is optimal.

General Algorithm To Build The Final Blocks

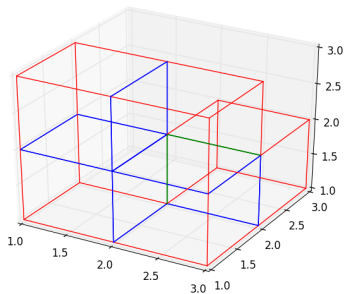
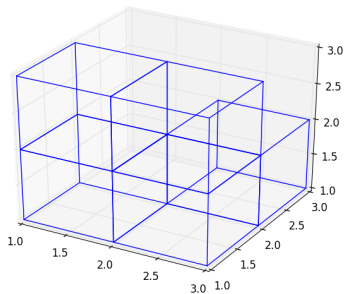
Algorithm 1 General Algorithm

```
1: mark all the blocks with the labels of the original block from which they were extracted
2: if all the blocks belong to the same original block then
3:   call function to get structured blocks
4: else if the original blocks are compatible (Call Routine to check for degeneracy) then
5:   merge all the blocks
6:   call function to build structured blocks
7: else
8:   divide the blocks in 3D representable sets (TODO)
9:   for all the 3D representable sets do
10:    call function to build structured blocks
11:   end for
12: end if

13: function BUILD_STRUCTURED_BLOCKS(3D representable blocks)
14:   call routine to detect concave edges (implemented)
15:   if there are no concave edges then
16:     all the blocks assigned to the processor form a structured block
17:     return the blocks structure
18:   else
19:     call routine to build structured blocks (TODO)
20:     return the blocks structure
21:   end if
22: end function
```

General Algorithm To Build The Final Blocks

Comparison between the original set of blocks and the final result where the concave edges are detected



Let L denote the original list of cuboids. Repeat the following steps as long as the list contains at least one mergeable cuboid.

- 1 Choose a mergeable cuboid (let it call A).
- 2 Among the cuboids that can be merged with A , select a cuboid B .
- 3 Merge cuboids A and B , yielding a new cuboid C .
- 4 Replace cuboids A and B by cuboid C in the list L .

- Implement refinement and heuristic algorithm for a more complicated structure.
- Compare with a deterministic optimization technique.
- Study various clustering methods for parallel optimization.



Thank You For Your Attention!